
Substitute Goods and Multi-Item Auctions

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What's New and What's Not

- ◆ Much of this presentation is a summary of known results about the important consequences for auctions of two conditions:
 - “goods are substitutes”
 - the coalitional value function is “bidder-submodular.”
- ◆ The relation between these two conditions is “known but opaque.”
- ◆ What is *new* is direct mathematical treatment that connects the two conditions tightly.

“Gross Substitutes”

SUBSTITUTE GOODS

Substitute Goods

- ◆ Each bidder n is described by a continuous, non-decreasing valuation function v^n that specifies the value to the bidder of any bundle x^n of goods.

- ◆ Our analysis makes use of Hotelling profit functions:

$$\pi^n(p) \equiv \max_{x^n \geq 0, x^n \in X} v^n(x^n) - p \cdot x^n$$

- ◆ By Hotelling's lemma, **goods are substitutes if and only if π is submodular.**

Implications for Equilibrium

- ◆ **Theorem.** If goods are (“strong”) substitutes for all bidders, then for every supply vector y there exists a price vector p such that y is in demand set at p .
- ◆ Applies to both divisible and indivisible goods.
 - Milgrom and Strulovici (2009): “strong substitutes”
 - Milgrom (2000)
 - Gul and Stacchetti (1999)
 - Kelso and Crawford (1982)

Aggregate Objective Function

- ◆ Let $v^S(y) \equiv \max_x \sum_{i \in S} v^i(x^i)$ subject to $\sum_{i \in S} x^i \leq y$,
 - Market demand can be computed by maximizing this function at any given prices. This works for both divisible and indivisible goods.
- ◆ Define $\pi^S(p) \equiv \sum_{i \in S} \pi^i(p)$.

Equilibrium Existence \Rightarrow Duality

- ◆ **Theorem.** For the discrete case as well as the concave case, for all S, y , $v^S(y) = \min_{p \geq 0} \sum_{i \in S} \pi^i(p) + p \cdot y$.
- ◆ **Proof.** Fix y and let \bar{p} be the corresponding market clearing price vector. Then, $y \in \operatorname{argmax}_z v^S(z) - \bar{p} \cdot z$. Hence, $\pi^S(\bar{p}) = v^S(y) - \bar{p} \cdot y$, so $v^S(y) \geq \min_{p \geq 0} \pi^S(p) + p \cdot y$.
By definition of π^S , for all p , $v^S(y) \leq \pi^S(p) + p \cdot y$. So, $v^S(y) = \min_p \pi^S(p) + p \cdot y = \min_p \sum_{i \in S} \pi^i(p) + p \cdot y$. **QED**

This treatment based on Ausubel and Milgrom (2002)

CORE AND BIDDER SUBMODULARITY

The Core

- ◆ Informally, the core consists of the “*competitive*” outcomes of an auction – the ones in which no group of bidders is willing to pay more to the seller for the items on offer.
- ◆ A feasible, individually rational auction outcome is *blocked* and hence *not* in the core if there is some coalition (necessarily including the seller) that could deviate to do better.
 - The seller would be paid more.
 - The deviating bidders would get better deals.

Maximal Core Payoffs

- ◆ Ausubel-Milgrom and Bikchandani-Ostroy.
 - Player 0 is the seller
- ◆ Theorem. For any set of bidders N , bidder i 's Vickrey payoff is her maximum payoff over all core imputations:

$$w(N) - w(N - i) = \max \{ \pi^i : \pi \in \text{Core}(N, w) \}$$

Proof Sketch

- ◆ Any imputation with a higher payoff for i is blocked by the coalition M_i .
- ◆ The imputation where i gets $w(N) - w(M_i)$, other bidders get 0, and the seller gets the rest is in the core.

Bidder Submodularity

- ◆ For any set of bidders S , the value of that coalition alone is zero, but the value if the seller with goods endowment \bar{x} is also included is:

$$w(S) = \max_x \sum_{j \in S} v^j(x^j)$$

subject to $\sum_{j \in S} x^j \leq \bar{x}$

- ◆ “Bidder submodularity” means that the function w is submodular, that is,

$$(\forall S, T \subset N) \quad w(S) + w(T) \geq w(S \cap T) + w(S \cup T)$$

Two Equivalent Formulations

- ◆ w is bidder submodular if and only if

$$(\forall S, i, j \notin S)$$

$w(S + i) - w(S) \geq w(S + i + j) - w(S + j)$ or equivalently,

$$w(S + i + j) - w(S) \geq (w(S + i + j) - w(S + j)) + (w(S + j + i) - w(S + i))$$

- ◆ These say roughly that
 - Expanding the set of bidders (by adding j) reduces the marginal contribution of each existing bidder i
 - The marginal contribution of the pair ij to the coalition $S+i+j$ is greater than the sum of their individual contributions

Vickrey Monotonicity

- ◆ In a Vickrey auction, each bidder's payoff is its marginal contributions to the coalition consisting of the seller and all other bidders: $w(S+j) - w(S)$.
- ◆ **Theorem.** For all sets of bidders S , each bidder's Vickrey payoffs falls weakly when an additional bidder enters the auction **if and only if w is bidder submodular.**

When is Vickrey in the Core?

- ◆ **Theorem.** The Vickrey auction outcome with bidders S is in $\text{Core}(S, w)$ for all $S \subseteq N$ **if and only if w is bidder submodular.**
- ◆ Informal arguments. In a Vickrey auction...
 - ... the outcome is efficient and, if w is bidder submodular, then each coalition of bidders $M \setminus S$ is paid no more than its incremental contribution, so the coalition S cannot block.
 - ...if w is NOT bidder submodular, then some pair of bidders $\{i, j\}$ gets MORE than its incremental contribution to some $S + i + j$, so coalition S can block.

Finally! What's New!

CONNECTING SUBSTITUTES TO BIDDER SUBMODULARITY

A Coalitional Value Calculation

- ◆ What is the value of a coalition consisting of a bidder with valuation v and another with an additive or linear valuation p ?

$$\begin{aligned}w(v, p) &= \max_{x \geq 0} v(x) + p \cdot (\bar{x} - x) \\ &= \pi(p) + p \cdot \bar{x}\end{aligned}$$

- ◆ *This ties dual functions directly to coalitional value functions.*

...and a second Calculation

- ◆ What is the value of a coalition consisting of a bidder with valuation v and two others with additive valuations p and p' ?

$$\begin{aligned}w(v, p, p') &= \max_{x \geq 0} v(x) + (p \vee p') \cdot (\bar{x} - x) \\ &= \pi(p \vee p') + (p \vee p') \cdot \bar{x}\end{aligned}$$

- ◆ *This ties a lattice notion in the dual directly to coalitional value functions.*

Substitutes and Coalitional Values

- ◆ Let V be the set of goods valuations from which bidder valuations are drawn.
 - Let V_{add} be the set of additive goods valuations.
 - Let V_{sub} be the substitutable valuations.
- ◆ Given a profile of valuations $v \in V^N$, let w_v be the corresponding coalition value function.
- ◆ **Theorem.** (Substitutes & Bidder-Submodularity)
Suppose that $V_{\text{add}} \subseteq V$. Then, V includes **only substitutes goods valuations** ($V \subseteq V_{\text{sub}}$) *if and only if* for all $v \in V^N$, w_v is **bidder-submodular**.

“Direct” Proof Ideas

- ◆ The definition of the coalitional value function is naturally given in terms of a maximization.
- ◆ Duality allows the maximization to be expressed as a dual minimization.
- ◆ Submodularity of the dual profit functions easily implies submodularity of the coalitional value functions.
- ◆ Hotelling’s lemma connects submodularity of the dual profit function to substitutes in demand.

Proof: Direct

- ◆ Suppose that $V \subseteq V_{sub}$. Think of a set S as represented by an indicator vector of 0's and 1's. Then for all $v \in V^N$,

$$\begin{aligned}w(S) &= \max_{x \geq 0} \sum_{i \in S} v^i(x^i) \text{ subject to } \sum_{i \in S} x^i \leq \bar{x} \\ &= \min_{p \geq 0} \max_x \sum_{i \in S} v^i(x^i) - p \cdot \left(\sum_{i \in S} x^i - \bar{x} \right) \\ &= \min_{p \geq 0} \sum_{i \in S} \pi^i(p) + p \cdot \bar{x} \\ &= \min_{p \geq 0} \sum_{i \in N} S^i \pi^i(p) + p \cdot \bar{x} \\ &= \min_{p \geq 0} F(p, S)\end{aligned}$$

- ◆ Since v is substitutable, π is decreasing and submodular, so F is submodular. Thus, w is submodular.

Proof: Converse

- ◆ Suppose the goods are not substitutes, so v is not submodular. Then there exist price vectors p and p' such that $\pi(p \wedge p') + \pi(p \vee p') > \pi(p') + \pi(p)$. But,

Coalition

Value

$$v, p \wedge p' \quad \pi(p \wedge p') + (p \wedge p') \cdot \bar{x}$$

$$v, p \wedge p', p \quad \pi(p) + p \cdot \bar{x}$$

$$v, p \wedge p', p' \quad \pi(p') + p' \cdot \bar{x}$$

$$v, p \wedge p', p, p' \quad \pi(p \vee p') + (p \vee p') \cdot \bar{x}$$

Recalling that $(p \vee p') \cdot \bar{x} + (p \wedge p') \cdot \bar{x} = p \cdot \bar{x} + p' \cdot \bar{x}$,

it follows that the coalitional value function is not submodular. **QED**

Vickrey & Substitutes Summation

◆ Theorem.

If $V_{\text{add}} \subseteq V$, then the five conditions below are equivalent.

1. $V \subseteq V_{\text{sub}}$
2. For all $v \in V^N$, $\pi_{w_v}^{\text{Vickrey}} \in \text{Core}(N, w_v)$
3. For all $v \in V^N$, $\pi_{w_v}^{n, \text{Vickrey}}(S)$ is non-increasing in S .
4. For all $v \in V^N$, losing bidders in the Vickrey auction have no profitable joint deviation.
5. For all $v \in V^N$, no Vickrey bidder can gain by the use of shills.

◆ Proof: See Milgrom (2004) chapter 8.

END